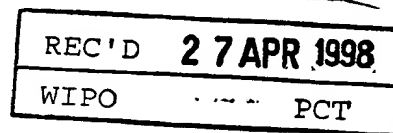


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Bescheinigung

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Patentanmeldung Nr. Patent application No. Demande de brevet n°

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PRIORITY DOCUMENT

Der Präsident des Europäischen Patentamts;
Im Auftrag

For the President of the European Patent Office

Le Président de l'Office européen des brevets
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Blatt 2 der Bescheinigung
Sheet 2 of the certificate
Page 2 de l'attestation

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Node supporting links having the ability to transfer longer messages than according to current MTP level 2

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Description

Node supporting links having the ability to transfer longer
5 messages than according to current MTP level 2

Background

Figure 1 shows the various protocol stacks for SS7
(Signalling system nr. 7) up to the MTP (message transfer
part) level. Five stacks are currently defined. The first one
10 is the well known stack for operation on 56/64kbit/s links.

Due to an increased bandwidth delay product the MTP (message
transfer part) level 2 (Q.703) is not ideally suited for
speeds significantly above 64kbit/s. The elements which are
problematic are window size, retransmission strategy, and the
15 error rate monitor.

Three different protocol stacks have been defined for use on
T1/E1 links (1.5/2 MBit/s) addressing some or all three of
these aspects.

The latest edition of Q.703 contains as a national option a
20 modification to the level 2 protocol which introduces 12 bit
sequence numbers and a different error rate monitor (second
column). Otherwise the procedures are not changed.

Recommendation Q.2119 defines frame-relay framing for SSCOP
(Service specific connection oriented protocol, Q.2110) to be
25 used on a raw E1/T1 link (third stack). Starting at SSCOP the
complete broadband protocol stack can thus be used on high
speed signalling links.

Finally Bellcore defines the complete ATM signalling protocol
stack starting at the ATM layer for use on T1 signalling
30 links, with certain restrictions in the ATM layer, like not
allowing multiple VCs (virtual channels) on a T1 link (column
4).

Lastly, the full ATM signalling protocol stack (column 5) could also be used in narrowband networks.

Besides the potentially vastly different link speeds which, however, pose no new interworking problems, the main difference between MTP level 2 based and SSCOP based signalling lies in the different maximum MSU length supported.

Of course there is no need to actually make use of the longer MSU length supported by the ATM links in an enhanced narrowband signalling network. Indeed the existing narrowband SS7 user parts would not even make use of the longer MSU length. We note, however, that the users of the SCCP can indeed generate messages in excess of 255 octets (the maximum data size supportable in single messages of the pre-96/97 SCCP). Such messages will be segmented before being delivered to the MTP. If such traffic would go via ATM links, avoiding the segmentation would benefit performance significantly.

Therefore we have the situation that use of the larger MSU sizes - where needed and possible - would be an additional welcome benefit of using the enhanced linksets.

SS7 Routing

25

Each node in an MTP network is identified by one signalling point code.

An MTP network is identified by the so-called *network indicator* in an MTP message.

30 Routing in the MTP is based on the so-called *destination (signalling) point code (DPC)* which identifies the destination of a *message signalling unit (MSU)* in an MTP network. In addition, the *signalling link selection field (SLS)* can be used to select between available routes of equal

priority (*combined linksets*) and to select a specific link within a linkset (a collection of links directly connecting two signalling points). No other information (like origination, MTP user, or MSU length) is generally evaluated
 5 for routing in the MTP.

The SCCP augments the MTP routing by providing additional functions to route on a so-called *global title* (GT), which can e.g. be a subscriber number of an 800-number. An SCCP routing on GT performs a process called *global title translation* (GTT) which derives the DPC of the final
 10 destination or the DPC of the next node (intermediate translator node) where the GT is further analyzed, eventually leading to the DPC of the final destination.

In addition to the GT the SCCP uses a so-called *subsystem number* (SSN) to identify the addressed SCCP user in the final
 15 destination.

This process also allows an SCCP message to cross MTP network boundaries.

In addition, the outcome of a GTT can depend on the
 20 availability status of the (next) destination. If the so-called *primary* destination, which would normally be the result of a GTT, or the addressed SSN is not available or reachable, an alternative destination can be the result of the GTT. This allows the SCCP to route messages to *backup*
 25 destinations (or backup intermediate translator nodes). Loadsharing between destinations is, in principle, also a possibility. Between two SCCP nodes the messages are routed by the MTP using the DPC provided by the SCCP.

30 State of the art

The interworking problem arising if use of longer messages in networks containing also linksets supporting only short messages is to be made has not been addressed in any detail.

Bellcore simply specifies that long messages destined for an MTP level 2 based link are to be discarded and that otherwise routing should be administrated accordingly.

5 A similar solution is proposed for the MTP based narrowband-broadband interworking in Q.2210.

For the SCCP the possibility is defined to convert long LUDT(S) messages into segmented short XUDT(S) messages.

10 All these solutions, however, require appropriate planning of the routes supporting the longer messages and/or will not make optimal use of the capabilities available.

An MTP level 3 protocol based approach to solve such problem is described in Q.701. This solution, however, is incomplete.

Addressing based solution

15 This invention proposes to use the addressing mechanisms provided in MTP and SCCP to solve or rather prevent the interworking problem. This works as follows:

Each node, which supports linksets having the ability to transfer longer messages than according to Q.703 (for example
20 SSCOP-linksets), is assigned a second point code (in addition to its *narrowband* point code), which will be called *broadband* pointcode, identifying its enhanced functions, i.e. those which can generate long messages. An example of such a network is given in figure 2. Routing tables in the MTP are
25 engineered so that these *broadband* signalling points are only connected via linksets supporting the longer message length (see tables 1 to 3 for an example). Non-enhanced nodes would have no knowledge about the broadband point codes in the MTP network (see table 5 for an example). For the interconnection
30 of the narrowband point codes and the non-enhanced nodes (i.e. the nodes having only *narrowband* point codes) all linksets, however, would be available.

Thus the nodes supporting the enhanced links (nodes identified also by the broadband signalling point codes) together with the enhanced linksets would form an overlay network which can transport longer messages (see figure 3).

- 5 Even nodes having only the enhanced linksets would be identified by a narrowband and a broadband point code.

It is, however, still possible for the SCCP to reach a node (having a narrowband and a broadband point code) to which no enhanced route is currently available by appropriately engineering the SCCP GT translation data if this should be
10 desired by the operator of the network.

GT translation in the SCCP of a node having a narrowband and a broadband point code is engineered so that physical destinations (intermediate translators or final destinations)
15 having a narrowband and a broadband point code have the broadband point code as the primary translation result and the narrowband point code as the backup translation result (see table 4).

As long as two signalling points are connected an enhanced
20 route will be used. If all enhanced routes between two nodes having a narrowband and a broadband point code fail communication between the nodes will be via the linksets supporting only short messages, using the narrowband point codes as addresses.

- 25 In addition, this solution can also be used for any new MTP users or appropriately modified existing MTP users like ISUP.

Similarly this solution is also suitable for interworking between narrowband and broadband signalling networks.

Note that an alternative solution would be to use a different
30 network indicator for the enhanced part of the signalling network which would have the advantage that there would be no restrictions in the available address space for point codes.

Table 1: MTP routing table in node a/A without link failure

destination	next node	
b	b	c
B	B	
c	c	b
d	b	c
D	B	

Table 2: MTP routing table in node a/A with link A to B failed
short messages can still reach all nodes via c

destination	next node	
b		c
B		
c	c	
d		c
D		

Table 3: MTP routing table in node a/A with link C-D failed
long messages to D not possible anymore

destination	next node	
b	b	c
B	B	
c	c	b
d	b	c
D		

**Table 4: SCCP global tile translation in node a/A for GT resulting in addressing the SCCP
 (or one of its users) in node d/D**

primary result (MTP address)	backup result (MTP address)
D (long message allowed)	d (segmentation required)

Table 5: MTP routing table in node c without link failure

destination	next node	
a	a	b
b	b	
d	d	b

Claims

1. Node supporting links having the ability to transfer longer messages than according to current MTP level 2, so-called enhanced links, to be identified by two signalling point codes one being used to identify/address functions/MTP users which can make full use of the longer message length, but both point codes being part of the same MTP network.

2. Node according to Claim 1, characterized by MTP routing tables supporting said enhanced links structured so that routing between the point codes used to identify/address functions/MTP users which can make full use of the longer message length uses only enhanced links.

3. Node according to Claim 1 or 2, characterized by SCCP translation functions supporting said enhanced links engineered that primary translation is to logical destinations reachable via said enhanced links and backup translation is to logical destinations reachable (also) via links based on MTP level 2 if translation results in a physical destination located in a node supporting said enhanced links.

4. Node supporting links having the ability to transfer longer messages than according to current MTP level 2, so-called enhanced links, to be identified by two signalling point codes one being used to identify/address functions/MTP users which can make full use of the longer message length, with the two point codes being part of different MTP networks, i.e. with a different network indicator, but all point codes being used to identify/address functions/MTP

users which can make full use of the longer message length located in the same MTP network.

5. according to Claim 4, characterized by

- 5 MTP routing tables supporting said enhanced links structured so that routing between the point codes used to identify/address functions/MTP users which can make full use of the longer message length uses only enhanced links.

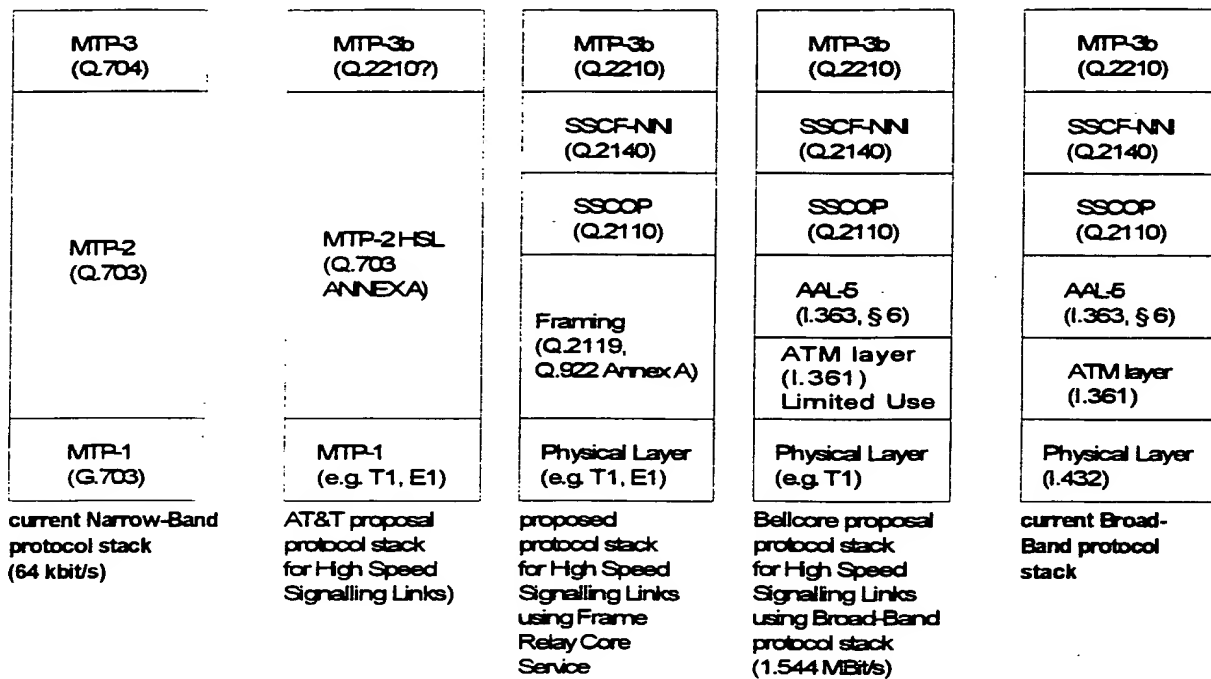
10 6. according to Claim 4 or 5, characterized by

- SCCP translation functions supporting said enhanced links engineered that primary translation is to logical destinations reachable via said enhanced links and backup translation is to logical destinations reachable (also) via
15 links based on MTP level 2 if translation results in a physical destination located in a node supporting said enhanced links.

20

25

Protocol Stacks proposed for High Speed Signalling Links

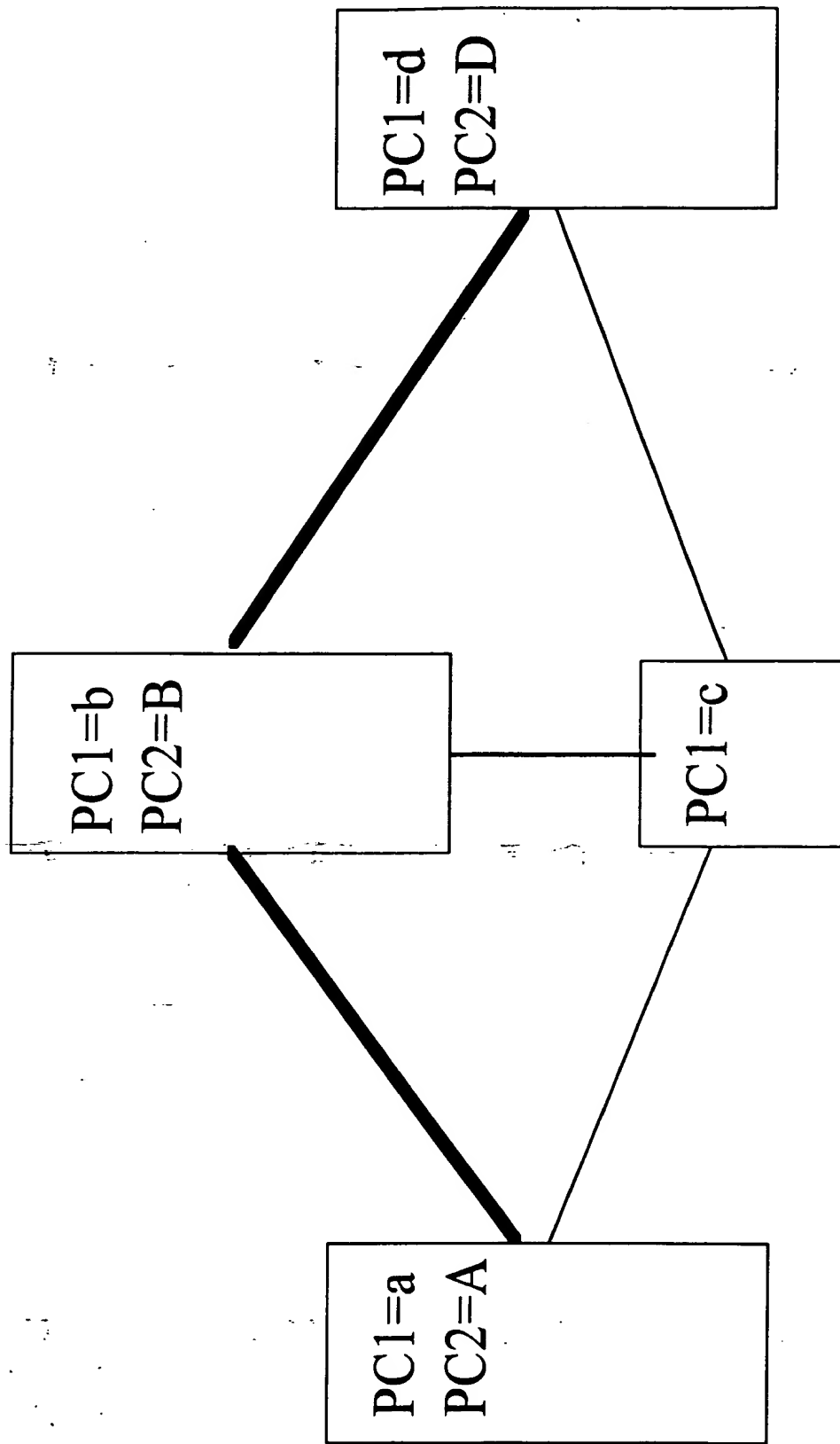


SSCF-NN: Service Specific Coordination Function at the Network Node Interface

SSCOP: Service Specific Connection Oriented Protocol

FIG 1

Figure 2 - physical network configuration



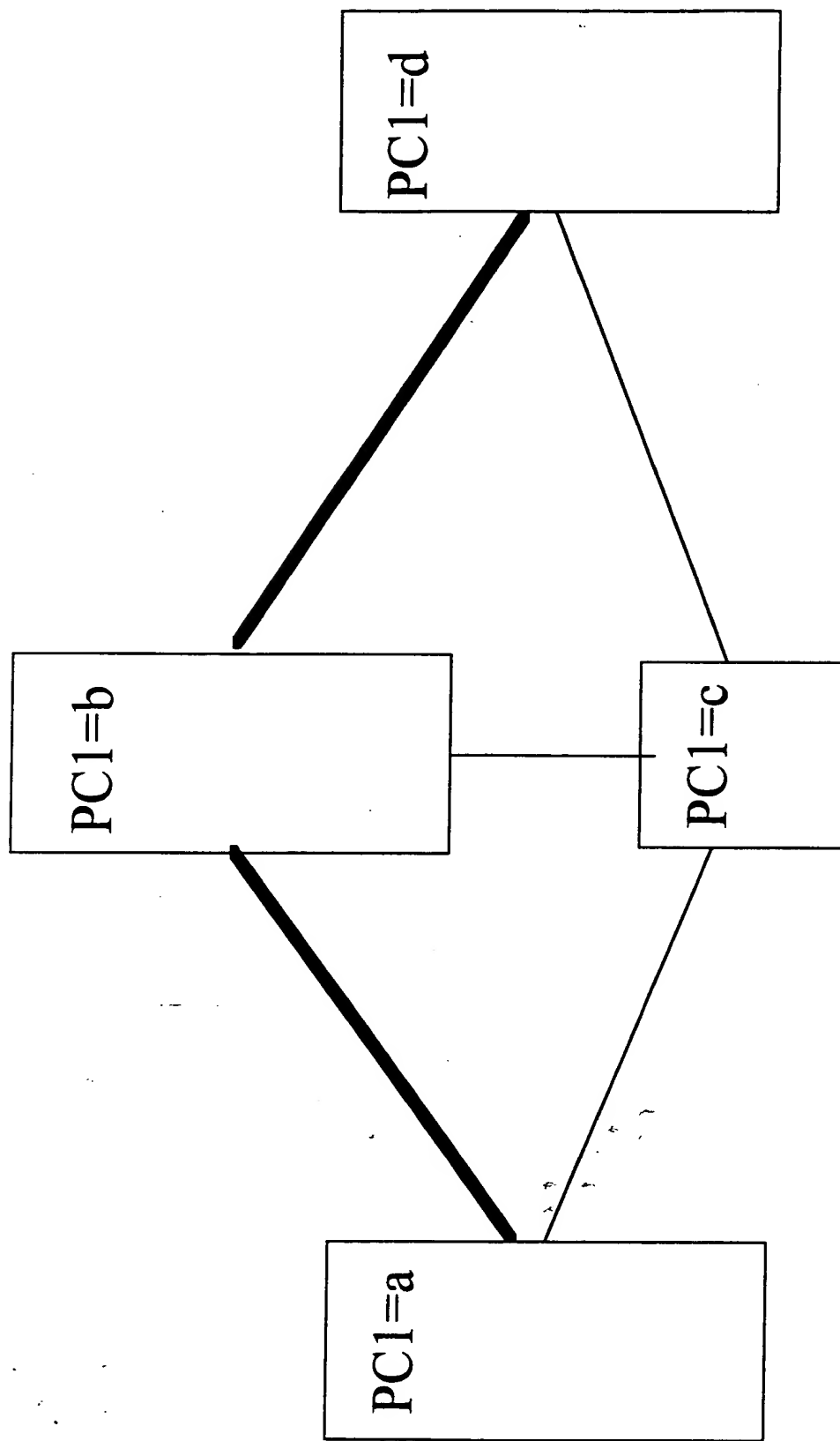
PC1 point code 1 (*narrowband point code*)

PC2 point code 2 (*broadband point code*)

— link/linkset based on SSCOP (Q.2110)

— link/linkset based on MTP level 2 (Q.703)

Figure 3 - logical network for short messages

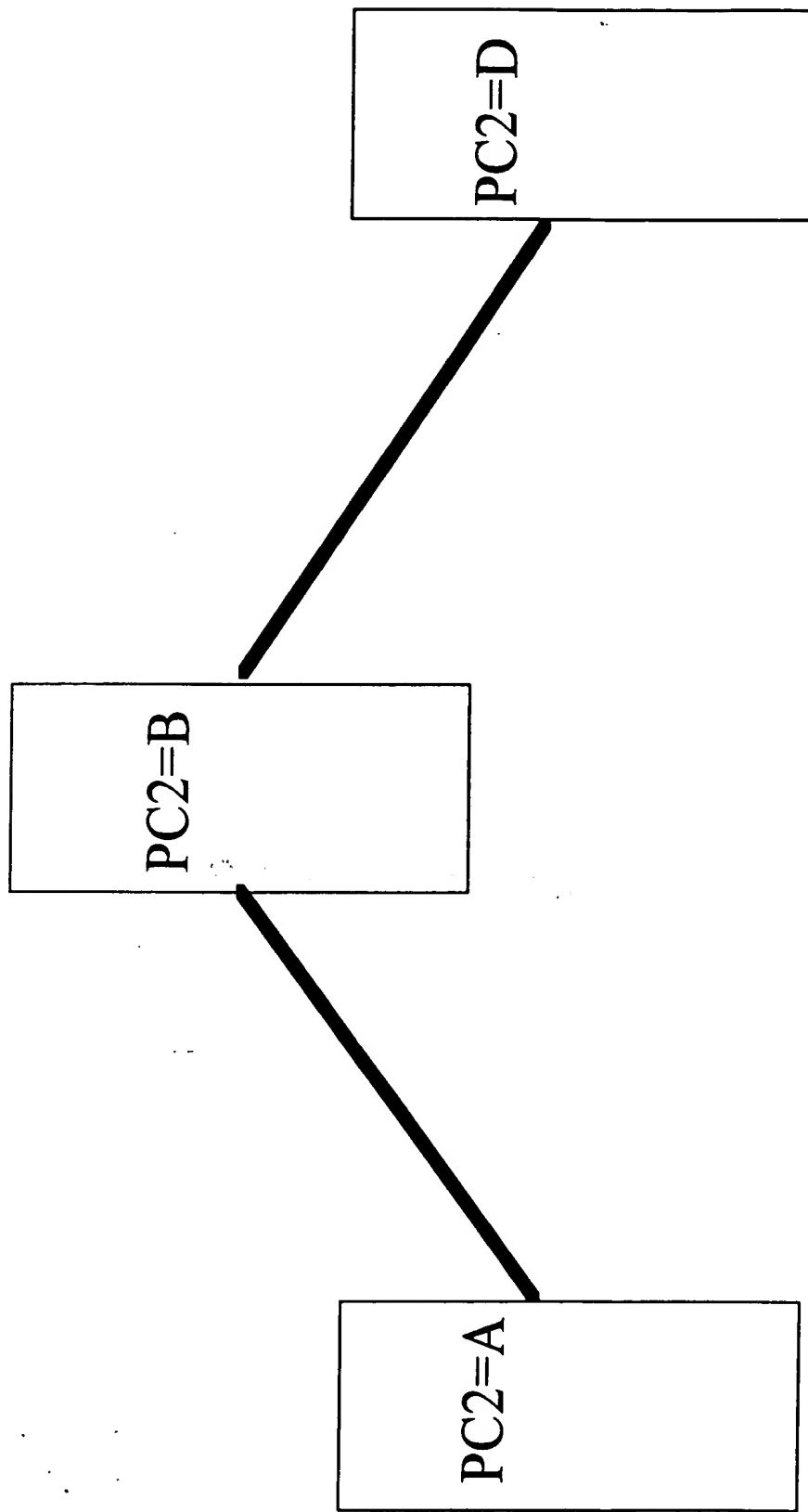


PC1 point code 1 (*narrowband point code*)

— link/linkset based on SSCOP (Q.2110)

— link/linkset based on MTP level 2 (Q.703)

Figure 4 - logical network supporting long messages



PC2 point code 2 (*broadband point code*)
— link/linkset based on SSCOP (Q.2110)

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